ForestDB: A Fast Key-Value Storage System for Variable-Length Key Strings

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B+Tree Limitations

- Not suitable to index variable-length long keys
  - Significant space overhead as entire key strings are indexed in non-leaf nodes
- Tree depth grows quickly as more data is loaded
- I/O performance is degraded significantly as the data size gets bigger
- Several variants of B+Tree were proposed
  - LevelDB (Google)
  - RocksDB (Facebook)
  - TokuDB (Tokutek)
  - WiredTiger (MongoDB)
ForestDB

- Key-Value storage engine developed by Couchbase Caching / Storage team

- Its main index structure is built from Hierarchical B+-Tree based Trie or HB+-Trie

- Significantly better read and write performance with less storage overhead

- Support various server OSs (Centos, Ubuntu, Debian, Mac OS x, Windows) and mobile OSs (iOS, Android)

- ForestDB paper published in *IEEE Trans. On Computers*
HB+Trie (Hierarchical B+Tree based Trie)

- Trie (prefix tree) whose node is B+Tree
  - A key is split into the list of fixed-size chunks (sub-string of the key)

Variable length key: `a83jgls83jgo29a...`  
Fixed size (e.g. 4-byte)

Search using Chunk1

Search using Chunk2

Search using Chunk3

Lexicographical ordered traversal

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Benefits

- When keys are sufficiently long & uniform random (e.g., UUID or hash value)
- When keys have common prefixes (e.g., secondary index keys)

Example: Chunk size = 4 bytes

Majority of keys can be indexed by first chunk
- There will be only one B+Tree on HB+Trie
- We don’t need to store & compare entire key strings